

Report on
CSFW-FCC-ICALP

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CSFW-FCC-ICALP

- 2006 July: Italy - Venice - S. Servolo Island
- CSFW: 5-7 July
 - FCC: 9 July
 - ICALP: 10-14 July



S. Servolo Island

General of FCC

- FCC(2nd): 1st: all accepted, 2nd: 10/13 accepted

- 5 minutes talk (12 talks: advertisement, ongoing)

- panel discussion on “non-determinism”:

Canetti, Micciancio, Mitchell, Pfitzmann, Palamidessi*, Segala*

(*prob. concurrent system)

next meeting

- CSFW: same place, same time

- FCC: ? join with ICALP/CSFW

- ICALP: Poland

FCC

(1) *Computationally Sound Secrecy Proofs by Mechanized Flow Analysis.*

Backes, Laud.

symbolic
foundation

(2) *Computationally Sound Compositional Logic for Security Protocols.*

Datta, Derek, Mitchell, Roy, Shmatikov, Turuani, Warinschi

symbolic
foundation

(3) *Language Design for Computationally Sound Communications Abstractions.*

Adao, Fournet.

symbolic
foundation

(4) *Soundness of Symbolic Equivalence for Modular Exponentiation.*

Lakhnech, Mazare, Warinschi.

positive
result

(5) *Sound and Complete Computational Interpretation of Symbolic Hashes in the Standard Model.*

Garcia, van Rossum.

positive
result

(6) *Soundness Limits of Dolev-Yao Models.*

Backes, Pfizmann, Waidner.

negative
result

(7) *Using Task-Structured Probabilistic I/O Automata to Analyze Cryptographic Protocols.*

Canetti, Cheung, Kaynar, Liskov, Lynch, Pereira, Segala.

real/ideal
foundation

(8) *Games and the Impossibility of Realizable Ideal Functionality.*

Backes, Datta, Derek, Mitchell, Ramanathan, Scedrov.

negative
result

(9) *An example of proving UC-realization with formal methods.*

Andova, Gjøsteen, Kråkmø, Mjøl̄snes, Radomirović.

symbolic
foundation

- tool based on BPW
- transform a protocol to a set of constraint
- abstract analysis deduce cryptographic secrecy?

related papers

Backes-ERORICS-2004, Backes-Pfitzmann-FST-TCS-2003

(2) *Computationally Sound Compositional Logic for Security Protocols.*

Datta, Derek, Mitchell, Roy, Shmatikov, Turuani, Warinschi.

symbolic

the state of art of
protocol composition logic by Mitchell et al

- history
- DH predicate and axioms for key exchange
- application: ISO-9798-3 key exchange
Kerberos V5

- a language based on process algebra
- no cryptographic primitive (as spi-cal)
(abstract secure primitives on securacy, authentication)

presented in ICALP 2006

complicated (by Mitchell)

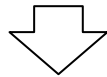
(4) Soundness of Symbolic Equivalence for Modular Exponentiation.

Lakhnech, Mazare, Warinschi.

positive

Extension of Abadi-Rogaway logic

($m \stackrel{\approx}{=} n$ in DY \rightarrow IND of m and n : encryption scheme: type0)



(..... : encryption scheme: IND-CPA
exponentiation : DDDH)

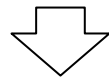
- **DynamicDDH**: (1) $\exp(x_1x_2-x_2x_3)$ (2) (3) neg dist from random
- slight modification of DY model
 - adv's knowledge
 - pattern
 - renaming (linear dep pres bij)

application: Burmester-Desmedt protocol

related work: D. Boneh?

Extension of Abadi-Rogaway logic

($m \cong n$ in DY \rightarrow IND of m and n : encryption scheme: type0)



(..... : encryption scheme: type
hash: Canetti's oracle hash)

- Canetti's oracle hash (1) $P[D_{\eta}(H(1^{\eta},x)=1)] - P[D_{\eta}(H(1^{\eta},y)=1)] < 1/p(\eta)$
(2) collision resistance
- slight modification of DY model
 - pattern
- sound & complete

(6) *Soundness Limits of Dolev-Yao Models.*

Backes, Pfitzmann, Waidner

negative

(1) protocols with hash prevent sound DY-model in BPW
(otherwise DY-model reverses hash)

(2) protocols with XOR prevent sound DY-model in BPW
(otherwise DY-model realizes signature creation/verif)
- restrict protocols allows sound DY-model

Related papers:

(1) “Limits of the Cryptographic Realization of Dolev-Yao-style XOR”
Backes, Pfitzmann (2005)

(2) “Limits of the Reactive Simulatability/UC of Dolev-Yao Models with Hashes”
Backes, Pfitzmann, Waidner (2006)

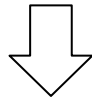
(7) *Using Task-Structured Probabilistic I/O Automata to Analyze Cryptographic Protocols*. Canetti, Cheung, Kaynar, Liskov, Lynch, **Pereira**, Segala.

real/ideal

extension of I/O Automata

- execution \rightarrow probabilistic execution [task scheduler]
- time bound (polynomial)

- prob. ver. of implementation (trace inclusion)
- prob. ver of simulation relation



accurate formalization of
IND between real and ideal in UC framework

Shoup's sequence of games

(8) *Games and the Impossibility of Realizable Ideal Functionality.*

Backes, Datta, Derek, Mitchell, Ramanathan, Scedrov.

negative

Scedrov → Mitchell → Derek → Backes

question: game description \Rightarrow real-ideal description

answer: generally negative

examples:

- multiparty coin-tossing
- bit-commitment (explained)
- shared random sequences

proof:

- create a game using ideal functionality.
- ideal realization \rightarrow information theoretic contradiction
the game

(9) *An example of proving UC-realization with formal methods.*
Andova, Gjøsteen, Kråkmo, Mjølunes, Radomirović.

symbolic

Proving universally composable theorem of
 (F_{PKI}, F_{SC}) - hybrid system for F_{SM}
by formal method

similar approach with Mitchell et al?

CSFW

many talks on foundation of access control/ authorization /policy

(1) *Cryptographically Sound Theorem Proving*: Mukhamedov, Ryan

- flaw against verified protocol
- manual proof -> fix and model check -> flaw in umbd num of participants

(2) *Refuting Claimed Security Proofs for Tripartite Key Exchange with Model Checker*: Choo

- finding new flaws
- tools based on AI planning system

(3) *Cryptographically Sound Theorem Proving*: Sprenger, Basin, BPW

- verification tool for BPW
- transform BPW representation for Isabelle